

The Best Structure of Mesh Cyphering Provisional Netting

¹Y.Jyothi ²Mr.D.Sattibabu

¹(M.Tech Student, Department of CSE, Godavari Institute of Engineering and Technology, Village Chaitanya Knowledge, Mandal Rajahmundry, Dist East godavari, A.P, India)

²(Associate Professor, Department of CSE, Godavari Institute of Engineering and Technology, Village Chaitanya Knowledge, Mandal Rajahmundry, Dist East godavari, A.P, India.)

Abstract : Network coding, the approach of operating arts coding operations on the contents of packets whereas in transit through the network, was originally developed for wired networks; recently, however, it's been additionally applied with success additionally to wireless ad hoc networks. In case, it's been shown that network coding will return significant performance gains, e.g., reduced energy consumption, in ad hoc networks. In this paper we propose and analyze exploitation linear programming, the better throughput that a multicast application are able to do with and while not network coding in unreliable ad hoc networks; we show that network coding achieves 65th higher throughput than typical multicast in a typical ad hoc network situation. The superiority of network coding, already established by the analytic results, is confirmed by simulation experiments.

Keywords - Network, Cyphering, Structure

I. INTRODUCTION

Recent results on the benefits of network coding in wired networks have excited a lot of interest within the subject and specifically, within the application of network coding to wireless ad hoc networks. Network coding refers to the fundamental notion of performing coding operations on the contents of packets throughout a network, and is mostly attributed to Ahlswede et al, who showed the service of the network coding for multicast in wired networks. The work of Ahlswede et al. was followed by different work by Koetter and M'edard who showed that codes with a straightforward, linear structure were enough to realize the capability of multicast connections in lossless, wire line networks. This result was augmented by ho et al, who showed that, in fact, a random construction of the linear codes was enough; The service of such random linear codes for reliable connection over lossy packet networks like wireless ad hoc networks was soon accomplished. In, a prescription for the economical operation of ad hoc networks is given, that proposes exploitation the random linear coding scheme of including optimization strategies for selecting the days and locations for injecting coded packets into the network. In this paper we deal with the issue of differentiating the most end-to-end throughput that a multicast connection can do with network coding given an unreliable ad hoc network. We offer 2 mathematical optimization formulations for max through multicast: one with network coding and one without; then, we analyze the most throughput that network coding achieves to the most throughputs that conventional multicast achieves in associate example network topology. We develop formulations supported the convex programming formulation of the minimum value multicast problem for network committal to writing given. In distinction to the minimum value multicast issues thought of, modeling the wireless medium rivalry constraints is crucial within the maximum throughput drawback that we consider. The mathematical programming formulations of the most throughput multicast drawback presented during this paper includes the wireless medium contention constraints. We use a technique similar to those proposed to model such constraints. In Network coding wired networks be initially estimate as a shared system for conclude the multicast theoretical (max-flow min cut) capability in wired networks. In wired multicasting, data is distributed from a group of supply nodes to a group of destination nodes over a multihop network wherever the intermediate nodes simply forward their received packets via a pre-determined look-up table (routing). Ahlswede et.al instructed the innovative notion of writing on layer-3 packets rather than look-up forwarding on specific outgoing links, and showed that network output may be increased. In a very network using NC, routers perform (random) linear combination of incoming layer-3 packets and broadcast the result to any or all its neighbors. Randomized linear network coding schemes were shown to be enough in achieving the data theoretical max-flow, min-cut sure on network capability. Necessary & enough condition by Koetter et al. whereas the idea of NC was developed for the network (IP) layer, it's typically been enforced in observe at higher layers, like the transport or application layers. An elementary reason on why network writing is helpful is predicated on the premise of synchronous transmission from many (source) nodes to one (receive) node. Whereas this can be possible in a much wired network whereby simultaneous transmissions are deemed „orthogonal“, a multihop wireless network is sort of totally different. Wireless could be a shared medium (at least for nodes among a typical transmission range) and there's no natural spatial orthogonality. therefore wireless multihop networks have relied on alternative varieties of

orthogonality- in time (TDMA) or frequency(FDMA)- to realize interference-free transmission Wireless Network writing (WNC) uses non-orthogonal transmission that, however, enable recovery of multiple packets to reinforce combination network output. Network coding in wireless networks the printed nature of wireless (coupled with network topology) determines the character of interference. Synchronous transmissions in a very wireless network generally end in all of the packets being lost (i.e. collision). A wireless network so needs a hardware (as a part of the mack functionality) to reduce such interference. Therefore any gains from network writing are powerfully wedged by the underlying hardware and can deviate from the gains seen in wired networks. Further, wireless links are generally half-duplex attributable to hardware constraints, i.e. a node cannot at the same time transmit and receive attributable to the shortage of enough isolation between the 2 methods. One in every of the potential application of WNC is in multicasting. A redistributed formulation to output optimization for the multicasting drawback was introduced. However, if further objectives like increasing output subject to delay constraints are thought-about, and then network codes should be together designed with mack. Authors qualify the impact of random access mack schemes (such as CSMA/CA) on performance of Old North State in all-to-all knowledge dissemination system. Epidemic algorithmic program for Rumour Spreading This work focuses on networks represented as graphs & distributed algorithms, wherever equally to our case, nodes don't have data regarding the nodes they're act with. At every spherical, every node indiscriminately chooses a communication partner among the nodes that are connected to that through a grip, and either "pushes" or "pulls" data. Broadcasting in Radio Communication Networks during this body of labor the wireless atmosphere is shapely as a graph, where, once a node transmits a message, it's received by all its neighbors is sending. Once more transmissions are divided into rounds, wherever in every spherical a set of the nodes transmits, in manner scheduled to reduce conflicts & maximize data spreading. The goal is to spread the data within the smallest range of rounds. Each centralized & redistributed algorithms are given. Indicative results include that, the matter is NP-hard, there exist. Static networks wherever the amount of needed rounds in $\Theta(\log^2 n)$. Whereas there exist mobile networks wherever the amount of needed rounds in $\Omega(n)$. Employing a similar model, the matter of minimizing energy consumption over a static wireless network was recently studied.

II. LITERATURE SURVEY

Multicasting exploitation NC may be divided into 2 tasks are Routing: Finding minimum-cost sub graphs to support multicast connections i.e., determinant the speed at that to inject coded packets on every arc and Coding: determinant the contents of those packets. Multicasting in Wired Networks Lun et al. given redistributed algorithms that calculate minimum value sub graphs for establishing single static multicast connections in wired & wireless networks that use writing. These algorithms, in addition to existing redistributed schemes for constructing network nodes represent a completely redistributed approach for achieving minimum-cost multicast. They observed that multiple synchronous multicast connections may be treated one by one, that is remarked as superposition writing however it's sub-optimal. Li & Li obtained, through applied math a necessary & enough condition for multicast rate feasibility & associate economical & distributed sub gradient algorithmic program for computing the utmost multicast rate. They finished that Old North State might not be instrumental in achieving higher easy lay multicast rates in most cases. Rather, it facilitates the look of considerably a lot of economical algorithms to realize such optimality. Chi.et al. projected that Old North State base routing algorithmic program for multicast capability. They showed that once the common node degree is high. The doable output of network writing based mostly Multicast is way over that of the shortest path distribution tree routing algorithmic program & slightly larger than that of the utmost rate distribution tree routing algorithmic program. Noguchi et al. projected a way for load equalization and observed that achieving easy lay flow victimization network writing will produce congestion. Li & Li showed that output improvement because of network writing in directed network is $O(|V|)$ and thus unbounded. For undirected networks with integral routing, there still exist configurations that are feasible with network coding but infeasible routing only for the multiple independent unicast transmissions. For undirected networks with integral routing, there still exist configurations that are feasible with NC but infeasible with routing only. Multicasting in Wireless Networks Xi & Yeh used NC to achieve minimum cost multicast in interference-limited wireless networks where link capacities are functions of the signal to interference-plus-noise ratio (SINR). They considered joint optimization of NC sub graphs with power control & congestion control without excessive control overhead & designed set of distributed, node-based scaled gradient projection algorithms and derived scaling matrices for fast, guaranteed global convergence. Ho et al. compared multicast network coding for a time varying wireless network model with interference-determined link capacities instead of collision based wireless model with fixed link capacities and showed that the gap in multicast capacity between NC & routing decreases relative to a collision-based wireless model with fixed-link capacities and the main advantage of NC is reduction in complexity of optimization and operation as NC significantly reduces complexity of dynamic back pressure algorithms used for optimization. In order to reduce cost and complexity

of coding, Zhang and Fan proposed to find nodes that need encoding instead of doing coding at all nodes. They used a modified Ford Fulkerson algorithm to obtain the maximum flow and encoding nodes in undirected graph instead of getting encoding nodes by sub tree decomposition as presented by Fragouli et al. Yuan et al. in proposed a general modeling and solution framework for the throughput optimization problem in wireless networks instead of cost optimization. In the framework, data routing, wireless medium contention and network coding are jointly considered to achieve the optimal network performance. The cross-layer optimization approach decay the initial problem into data routing sub-problems at the network layer, and power allocation sub-problems at the physical layer taking into account physical layer interference. Fragouli et al. proposed & argued that main benefits of network coding in a wireless environment might manifest in situations where topology dynamically changes and operation is restricted to distributed algorithms that do not employ knowledge about the network environment. They proved that network coding can attempt benefits of a factor of $\log n$ in terms of energy efficiency. Lun et al. showed how network coding, combined with distributed flow optimization, gives a practical approach for uni casting that promises to significantly outperform the present approach of end-to-end or link-by-link re-transmission combined with route optimization for any performance measure which increases with the number of transmissions made by each node. Katti et al. proposed COPE, architecture for wireless mesh networks. They addressed the common case of unicast traffic, dynamic and potentially bursty flows, and practical issues facing the integration of network coding in the current network stack. The testbed deployment results display that COPE largely development network throughput. The gains vary from a few percent to several folds depending on the traffic pattern, congestion level, and transport protocol. Due to coding, COPE has to send less number of packets and as a result load on bottleneck link reduces, giving it double advantage. Sengupta et al. obtained a theoretical formulation for computing the output of network writing on any wireless network topology and any pattern of simultaneous unicast traffic sessions. They advocated that routing be created alert to network coding opportunities instead of, as in COPE, being oblivious to that. They studied the trade-off between routing flows “close to every other” for utilizing writing opportunities and “away from each other” for avoiding wireless interference and given a way for computing the source-destination routes and utilizing the most effective writing opportunities from available ones so on maximize the output. Network writing to realize energy savings in a very wireless ad-hoc network, wherever every node of the network is that the supply sending data to any or all alternative nodes, Fragouli et al. projected totally distributed algorithm to perform network coding whereas addressing sensible problems like forwarding issue, managing generations and impact of transmission vary. Li et al. applied network writing to deterministic broadcast approaches, leading to vital reductions within the range of transmissions within the network. They proposed 2 algorithms, that rely only on native two-hop topology data and makes use of opportunist paying attention to cut back the amount of transmissions: 1) an easy XOR-based writing algorithm that gives gains up to forty fifth compared to a non-coding approach and 2) a Reed-Solomon based mostly writing algorithm that determines the optimal cryptography gain achievable for a coding algorithm that depends only on native data with gains up to 61. Gkantsidis and Rodriguez in used network coding for content distribution of large files. The randomization introduced by the coding process eases the scheduling of block propagation. They showed that expected file download time and robustness of the system is significantly improved with network coding. Hamra et al. showed that network coding can improve the performance of the file sharing application in wireless mesh networks but not as in wired networks inspite of wireless broadcast advantage and cited reasons for the behaviour. They also determine the main framework that influence the performance of network coding in wireless environment and showed how these parameters interact with each other and influence the behavior of network coding. They also identified the main parameters that influence the performance of network coding in wireless environment and showed how these parameters interact with each other and influence the behavior of network coding.

III. RELATED WORK

Several solutions are proposed with the aim to improve the efficiency of broadcasting in ad hoc networks. A widely applied approach is based on probabilistic message transmissions, within which the messages generated by a traffic supply are rerun by alternative nodes with an exact probability. This answer reduces the overhead with respect to easy flooding, but it may cause a low delivery ratio, particularly within the case of scarcely connected nodes. In counter-based schemes, instead, a node determines whether to rerun a packet by counting how many identical packets it receives over a given quantity. In distance-based broadcasting, nodes create use of their distance from the previous sender as call criteria. The most idea is to enable those nodes that are situated farther away from the sender to rebroadcast the packet, therefore on increase the message spatial progress. Broadcasting techniques are recently studied also for the dissemination of alarms or warnings in transport networks. In broadcast packets include data regarding the sender position and also the message-propagation direction: whenever a node receives a duplicate of a broadcast message, it rebroadcasts the packet as long as the node follows the sender. The answer presented, instead, doesn't need position-based information. As for broadcasting supported network cryptography, a sensible scheme is planned, wherever network

cryptography is applied only by a set of nodes that are chosen as forwarders; Also, through promiscuous mode, a forwarder will decide not to send a packet if it hears that every one of its neighbors have received it. Albeit the theme provides smart performance, it depends on the exchange of neighborhood data and on the partial dominant pruning algorithmic rule conferred, which can cause important overhead and add complexness relatively to probabilistic approaches. An irregular broadcast scheme supported network coding is planned. There, the network coding parameters are finely tuned for the case of a grid topology thus on reduces the packet loss probability with relation to easy flooding. However, such edges will be perceived solely in dense networks where the packet loss probability because of collisions will be on top of the probability of coding failure. The impact of many IEEE 802.11-based mac protocols on the performance of network cryptography for broadcast communications is analyzed. Although relevant, we highlight that to the simplest of our knowledge only many works exist on network cryptography for ad hoc networks, relating sensible scenarios or exploitation realistic network simulators. Thus, one of the most goals of our work is to assess the performance of network cryptography for message broadcasting by exploitation the network machine ns-2 as an analysis tool. Moreover, we implement an efficient network cryptography scheme that doesn't need nodes to have information on their two-hop neighbors; we compare its performance against each simple flooding and deferred broadcast, the latter together with the solution proposed.

IV. EXISTING SYSTEM

The capability of wired networks is improved by network coding (NC), which might totally utilize the network resources. Because of this advantage, a way to use NC in wireless ad hoc networks has been intensively studied in recent years. Necessary to design the NC in wireless accidental networks with interference to realize the development on system performance like sensible place and delay/good place tradeoff. A very important work by Liu et al in introduced the observation that only a continuing issue of output improvement is caused to k-dimensional random static networks. It had been indicated in their results that order improvement of output scaling laws is achieved by adopting RLNC in MANETs. The probability that the random linear NC was valid for a multicast connection drawback on an arbitrary network with independent sources; It had been indicated in their results that order improvement of throughput scaling laws is achieved by adopting RLNC in MANETs.

V. PROPOSED SYSTEM

We have analyzed the NC configuration in each static and mobile ad hoc network to optimize the delay/ smart place tradeoff and therefore the smart place with the thought of the throughput loss and coding loss of NC. The mobile model, the two-hop relay scheme and therefore the flooding scheme are proposed for each random independent and identically distributed (i.i.d.) quality model and random walk model with random linear NC. The output loss and coding loss of NC that are treated because the overhead of NC, are also considered; the coding loss is caused by coding failure of RLNC. Since the NC coefficients are indiscriminately elect from Galois field F_q , the destination might not decipher the k original packets with success. The scaling laws of NC overhead, that weren't considered in most previous works in wireless ad hoc networks. The theoretical results indicate that the good place and delay/good place trade-off cannot be improved so as sense by using NC once considering the throughput loss and coding loss.

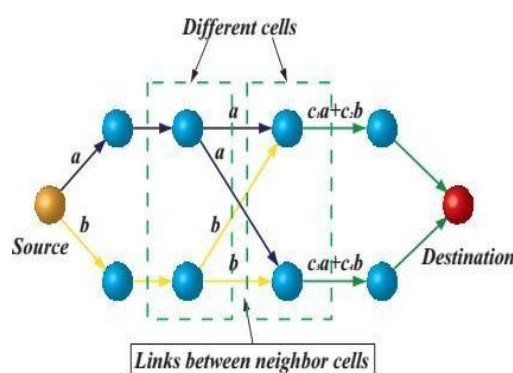


Fig.1. System Architecture

VI. FRAME WORK

ANALYSIS FOR STATIC NETWORKS

In this section, the static network model is introduced initially, which incorporates the configuration and transmission model. Moreover, we propose the transmission scheme for this model, and also the corresponding unit square area and delay are analyzed supported the thought of output loss and coding loss.

A. Network Topology

We specialize in the networks that contains n randomly and equally distributed static nodes in an exceedingly unit square area. These nodes are randomly sorted into source-destination (S-D) pairs.

B. Transmission Model

During this paper, we adopt the protocol model that could be a simplified version of physical model since it ignores the long distance interference and transmission. Moreover, it is indicated that the physical model may be treated because the protocol model on scaling laws once the transmission is allowed if the Signal to Interference Noise ratio (SINR) is larger than a given threshold.

C. Transmissions scheme for Static Networks

During this model, 3 sorts of nodes are concerned, i.e., supply node, relay node and destination node. Every node within the network could act joined or a number of the 3 roles. Moreover, the PNC is adopted during this theme since the configuration is mounted. This could be complete and calculated in an exceedingly computing center of this network. Since the network is therefore massive that every node hardly is aware of and stores all of the network writing coefficients, we only specialize in the case that every node doesn't recognize the network writing coefficients of others. Moreover, the case that every node is aware of the entire network writing coefficients is mentioned at the tip of this section. For the transmission theme, the k packets (as a generation group) are transmitted in an exceedingly digraph with the minimum total euclidean distance between connected nodes. within the static networks, the „when-to-stop“ signal is transmitted by handshake, and also the whole unicast session won't stop till $(1 + \epsilon)k$ completely different packets inward at destination

ANALYSIS FOR MOBILE NETWORKS

In this section, we'll introduce the mobile network models that embody configuration, quality models and transmission model. Moreover, the transmission schemes are planned for each of the quality models, and also the corresponding smart place and delay are derived within the similar method.

A. Network Topology

We have a tendency to study the mobile networks that contains n mobile nodes in an exceedingly unit sq. space. These nodes are randomly sorted into source-destination (S-D) pairs.

B. Mobility Models

In mobile networks, the overall region (the unit sq.) is split into $m = \Theta(n)$ square cells rather than $\Theta(\log n)$, wherever $m < n$. Our work may be applied to several quality models, and that we primarily target the random i.i.d. quality model and stochastic process model, that are outlined as follows. Random i.i.d. quality model: In random i.i.d. quality model, every node is in an exceedingly randomly designated cell severally and identically within the next time-slot, which suggests that every node is in any cell with an equivalent chance. As a result, the configuration changes drastically in each time-slot, and also the network behavior can't be foretold. This model is additionally adopted. stochastic process model: In stochastic process model, if a node is in cell $(i, j) \in \mathbb{Z}^2$ at the present time-slot, it'll be within the 5 cells: (i, j) , $(i + 1, j)$, $(i - 1, j)$, $(i, j + 1)$, $(i, j - 1)$ with an equivalent chance within the next time-slot. During this model, the speed of every node is restricted, so the model will therefore higher describe the realistic node quality than random i.i.d. quality model. Previous literatures adopting the stochastic process quality model;

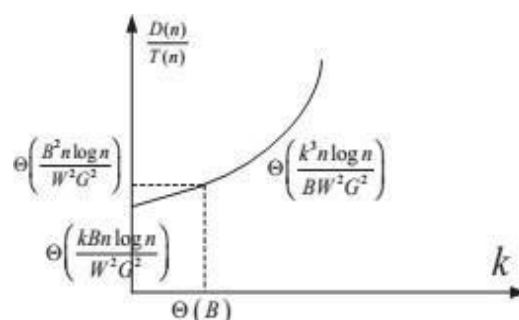
C. Transmission Schemes for Mobile Networks

The subsequent schemes area unit applicable to each random i.i.d. quality model and random walk model; firstly, we have a tendency to outline three sorts of transmissions: supply to relay (S-R), relay to relay (R-R) and relay to destination (R-D). Additionally, supply to destination transmission conjointly belongs to R-D. Once the relay receives a replacement packet, it combines the packet it's with the one it receives by randomly designated coefficients, so generates a replacement packet. Coinciding transmission in one cell isn't allowed since the receiver is difficult to get multiple csis from totally different transmitters at an equivalent time. Hence, we have a tendency to use the random linear network writing for mobile models. Specifically, we have a tendency to adopt 2 themes area unit 2-hop relay scheme, R-R transmission isn't allowed during this theme. The entire

packets are unit transmitted from supply to destination by at the most a pair of hops. The chance that either S-R or R-D is chosen is $1/2$. Every packet is deleted t_d seconds once its generation, wherever $t_d = \Omega(D(n))$ is determined supported the delay of the networks. Once S-R transmission is chosen, the seven supplies can randomly choose a node within the same cell and transmit a combined packet of k original packets with coefficients which are willy-nilly designated from F_q . Once R-D transmission is chosen, the relay can transmit the corresponding packet to the destination so delete this packet. The destination decodes the network writing once it receives $(1 + \epsilon)k$ completely different packets, wherever ϵ could be a positive constant. Flooding theme, all the 3 transmissions are allowed, and also the chance that one amongst them is chosen is $1/3$. The delete time t_d for this theme is additionally $t_d = \Omega(D(n))$. The R-D transmission is that the same as in 2-hop relay theme. Once S-R transmission is chosen, the supply can transmit a combined packet of k original packets with network writing coefficients to any or all the nodes among an equivalent cell. Once R-R transmission is chosen, one node is willy-nilly designated, and it'll transmit one amongst its packets to the opposite nodes within the same cell equiprobably. Once receiving the packet, every node during this cell combines the packet with an equivalent session packet it's (if there's one) as in (3), wherever $G = 1$. Afterwards, it'll store the mixture in its memory and delete the received packet in addition because the previous packet of this session. The coding time of destination is that the same as in 2-hop scheme.

VII. EXPERIMENTAL RESULTS

In this session, we discuss the results higher than and optimize the delay/goodput trade-off and goodput for each static and mobile network. The corresponding optimal information size B , generation size k and network committal to writing field F_q also are derived. Besides, we compare the results with no NC case. The information size for network while not NC. Since there's no network committal to writing during this case, the output of it will be treated because the goodput. In addition, it should be noted that the units for B , $T(n)$, $D(n)$ and also the tradeoff are bits, bits/sec, sec and sec²/bits, severally. The relations between the delay/goodput trade-off and k are illustrated for every model and scheme. Moreover, the enhancements of network coding;



VIII. CONCLUSION

In this paper, we've studied regarding the network coding and its advantages in wired and wireless networks. During this paper, we tend to analyze the network coding configuration in each static and mobile ad hoc networks therefore on optimize the delay/good put trade-off and also the good put with the thought of the throughput loss and decoding loss of network coding. We have additionally studied regarding the researches challenges in Network coding.

REFERENCES

- [1] T. Berger, "Multiterminal source coding," The Information Theory Approach to Communications, G. Longo, Ed. New York: Springer-Verlag, Aug. 1977.
- [2] R. Ahlswede, N. Cai, S. Li, and R. Yeung, "Network information flow," Information Theory, IEEE Transaction on, vol. 46, no. 4, pp. 1204-1216, Jul. 2000.
- [3] R. Ahlswede, N. Cai, S. Li, and R. Yeung, "Linear network coding," Information Theory, IEEE Transaction on, vol. 49, no. 2, pp. 371-381, Feb. 2003.
- [4] J. Liu, D. Goeckel, et al., "Bounds on the gain of network coding and broadcasting in wireless networks," Proc. of IEEE INFOCOM 2007, Anchorage, Alaska, USA, May 2007.
- [5] C. Zhang, X. Zhu and Y. Fang, "On the improvement of scaling laws for large-scale MANETs with network coding," Selected Areas in Communications, IEEE Journal on, vol. 27, no. 5, pp. 662-672, Jun. 2009.

- [6] J. Huang, V. Subramanian, R. Agrawal, and R. Berry, "Joint Scheduling and Resource Allocation in Uplink OFDM Systems for Broadband Wireless Access Networks," *IEEE Journal on Selected Areas in Communications*, vol. 27, no. 2, pp. 226 - 234, Feb. 2009.
- [7] J. Huang, V. Subramanian, R. Agrawal, and R. Berry, "Downlink Scheduling and Resource Allocation for OFDM Systems," *IEEE Transactions on Wireless Communications*, vol. 8, no. 1, pp. 288 - 296, Jan. 2009.
- [8] C. Zhang, Y. Fang and X. Zhu, "Throughput-delay tradeoffs in large-scale MANETs with network coding," *proc. of IEEE INFOCOM 2009*, Rio de Janeiro, Brazil, Apr. 2009.
- [9] Q. Yan, M. Li, Z. Yang, et al., "Throughput analysis of cooperative mobile content distribution in vehicular network using symbol level network coding," *Selected Areas in Communications, IEEE Journal on*, vol. 30, no. 2, pp. 484-492, Feb. 2012.
- [10] S. Karande, Z. Wang, H. Sadjadpour, J. Garcia-Luna-Aceves, "Multicast throughput order of network coding in wireless ad-hoc networks," *Communications, IEEE Transactions on*, vol. 59, no. 2, pp. 497-506, Feb. 2011.
- [11] K. Lu, S. Fu, Y. Qian, and H. Chen "On capacity of random wireless networks with physical-layer network coding," *Selected Areas in Communications, IEEE Journal on*, vol. 27, no. 5, pp. 763-772, Jun. 2009.
- [12] Z. Yang, M. Li, W. Lou, "Codeplay: live multimedia streaming in VANETs using symbol-level network coding," *Wireless Communications, IEEE Transactions on*, vol. 11, no. 8, pp. 3006-3013, Aug. 2012.
- [13] S. Sharma, Y. Shi, et al., "Is network coding always good for cooperative 13 communications?" *Proc. of IEEE INFOCOM 2010*, San Diego, CA, USA, Mar. 2010.
- [14] A. Asterjadhi, E. Fasolo, et al., "Toward network coding-based protocols for data broadcasting in wireless ad hoc networks," *Wireless Communications, IEEE Transactions on*, vol. 9, no. 2, pp. 662-673, Feb. 2010.
- [15] D. Niu and B. Li, "Analyzing the resilience-complexity tradeoff of network coding in dynamic P2P networks," *Parallel and Distributed Systems, IEEE Transactions on*, vol. 22, no. 11, pp. 1842-1850, Nov. 2011.